



# Optimizing the RJMS of an Academic HPC & Research Computing Facility

Dr. S. Varrette, Dr. E. Kieffer, Dr. F. Pinel

University of Luxembourg (UL), Luxembourg

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hpc.uni.lu

hpc@uni.lu

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# **Summary**

- Introduction: Context and Motivations
- 2 Toward Improved User Job Management through a Novel RJMS configuration
- 3 Impact Analysis of the Updated Models and Policies through Workload Performance Evaluation
- 4 Conclusion & Perspectives





Introduction: Context and Motivations

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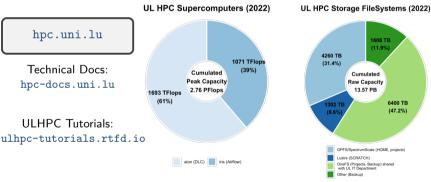
Technical Docs:

**ULHPC Tutorials:** 

Introduction: Context and Motivations

# **UL HPC Facility**

- Managed and operated since 2007 (Dr. S. Varrette & Co.)
  - → 2nd Largest HPC facility in Luxembourg after EuroHPC MeluXina





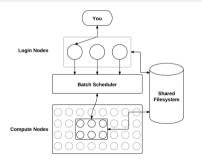


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# **Resource and Job Management Systems**

- Resource and Job Management System (RJMS)
  - → Glue for a parallel computer to execute parallel jobs
  - Goal: satisfy users demands for computation
    - √ assign resources to user jobs with an efficient manner







# **Resource and Job Management Systems**

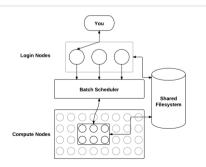
- Resource and Job Management System (RJMS)
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#### HPC Resources:

- → Nodes (typically a unique IP address)
  - √ Sockets / Cores / Hyperthreads
  - ✓ Memory
  - ✓ Interconnect/switch resources
- → Generic resources (e.g. GPUs)

#### Strategic Position

- → Direct/constant knowledge of resources
- $\hookrightarrow$  Launch and otherwise manage jobs





Introduction: Context and Motivations



https://slurm.schedmd.com/

#### Slurm on ULHPC clusters

- ULHPC uses Slurm for cluster/resource management and job scheduling
  - → Simple Linux Utility for Resource Management
  - → Reference RJMS serving most of Top500 systems
    - √ official documentation, tutorial, FAQ



Introduction: Context and Motivations

# SIUCM workload manager

#### Slurm on ULHPC clusters

- ULHPC uses Slurm for cluster/resource management and job scheduling

https://slurm.schedmd.com/

- Seminal configuration part of flagship iris production release (in 2017)
  - → migration from OAR RJMS
  - $\hookrightarrow$  Slurm provides superior scalability and performance [JSSPP12], inherent compatibility with multiple distributed libraries (Dask, IPyparallel...) and MPI suits (though PMI[x])
    - ✓ brings a more convenient and flexible interface for non-specialists
  - $\hookrightarrow$  inspired from other HPC centers (LLNL, Niflheim, CSCS...) & Simulators [PEARC18]

[JSSPP12] Y. Georgiou, M. Hautreux "Evaluating scalability and efficiency of the resource and job management system on large HPC clusters", in W. on Job Scheduling Strategies for Parallel Processing (LNCS JSSPP12), Springer, pp. 134–156 (2012).

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#### After 3 years of production on iris...

(1/2)

- Over-complexified setup for partition/QOS
  - → integration of hetegeneous hardware (GPU nodes, Large-memory) starting 2018





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  - - ✓ batch partition (Dual-CPU regular nodes) saturated and over-used
    - √ dedicated resources for interactive or long (regular nodes only) under-used





#### After 3 years of production on iris...

(1/2)

- Over-complexified setup for partition/QOS
  - → integration of hetegeneous hardware (GPU nodes, Large-memory) starting 2018
  - → under vs. over used partitions
    - ✓ batch partition (Dual-CPU regular nodes) saturated and over-used
    - √ dedicated resources for interactive or long (regular nodes only) under-used
  - → for each partition, associated QOS queue named qos-partitionname>[-XXX]
    - ✓ allows for specific research groups/industrial partners privileged/exclusive access
    - ✓ Ex: qos-batch (default), qos-batch-[001-...] (Group Prof. / Partner  $XX_{\{1-...\}}$ )
    - ✓ Ex: qos-gpu (default), qos-gpu-[001-...] (Group Prof. / Partner YY (1-...)
    - ✓ Ex: gos-covid ultra high priority jobs to support fight against COVID-19
  - → In practice, this setup brought frustration, jealousy & confusion within users





#### After 3 years of production on iris...

(2/2)

- No cross-partition QOS except best-effort (preemptible jobs):
  - → interactive jobs on non-regular nodes (GPU, large-memory) artificially complexified
  - → each specific usage treated by dedicated QOS (Ex: qos-batch-0\*, qos-covid)
    - ✓ no global priority level (low→urgent), unreadability of QOS objectives vs. partition





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- Fairsharing relying on the Depth-Oblivious Fair-share Factor algorithm
  - → variant of classical fair-share factor, increases usable priority ranges
    - √ very complex algorithm, hard to explain impacts/issues with job priority
  - → no policy for raw-share attribution lead to unfair situation
  - → hard to evaluate/size associated Slurm parameters

Ex: PriorityWeight\* (among others)





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Ex: PriorityWeight\* (among others)

- Incomplete Account hierarchy
  - $\hookrightarrow$  not able to cover new workload requests associated to funded projects/training events
  - $\hookrightarrow$  novel auditing capabilities on the platform usage / cost model requested
    - ✓ ... and need to be integrated in the Slurm configuration





Introduction: Context and Motivations

# Toward a Novel RJMS Setup

- Acquisition & integration of new liquid-cooled supercomputer aion
  - $\hookrightarrow$  occasion to deeply review and optimize the seminal configuration
  - $\hookrightarrow$  mitigate the identified pitfalls & take advantage of experience gained



# Toward a Novel RJMS Setup



- $\,\hookrightarrow\,$  occasion to deeply review and optimize the seminal configuration
- $\hookrightarrow$  mitigate the identified pitfalls & take advantage of experience gained

#### In this talk: ULHPC Slurm configuration 2.0

- Re-defining the partition, QOS and accounting model
  - → offer a more focused and easy-to-use configuration
- Rewiew fair-sharing model giving incentives to good practices
  - → take into account monetary contribution to increase priority
- rethinking global resources limits set to the partitions / accounts associations / QOS
- 4 consolidating the RJMS setup for HA services
- 6 define common [federated vs. multi-cluster] scheduling / accounting policy







#### Toward Improved User Job Management through a Novel RJMS configuration

# Summary

- Toward Improved User Job Management through a Novel RJMS configuration
- Impact Analysis of the Updated Models and Policies through Workload Performance Evaluation





#### **ULHPC Slurm Partitions 2.0**

-p, -partition=<partition>

- Tied to global types/classes of available computing nodes batch[,gpu][,bigmem]
- New floating partition interactive across all nodes for (short and quick) tests
  - → selection of the expected resource type through feature -c broadwell,skylake,gpu,volta[32]...
  - → backfill scheduling enabled and optimized to favor interactive and/or small jobs if queued
- Max wall clock time for user job on non-floating partition reduced from 5 to 2 days

AION Partition	Туре	#Node	PriorityTier	DefaultTime	MaxTime	MaxNodes
all (sysadmins)	hidden	318	100	10h	5d	UNLIMITED
interactive	floating	318	100	30min	2h	2
batch		318	1	2h	48h	64

IRIS Partition	Туре	#Node	PriorityTier	DefaultTime	MaxTime	MaxNode	es
 all (sysadmins)	hidden	196	100	10h	5d	UNLIMIT	ED
interactive	floating	196	100	30min	2h	2	
batch		168	1	2h	48h	64	
gpu		24	1	2h	48h	4	881
bigmem		4	1	2h	48h	1	UNIV



### **ULHPC Slurm QOS 2.0**

--qos=<qos>

- ullet New Cross-partition QOS, mainly tied to priority level (low o urgent)
  - $\hookrightarrow$  Simpler names than before (i.e. no more qos- prefix)
  - $\hookrightarrow$  special **preemptible QOS** kept for best-effort jobs: besteffort
  - → new long QOS allows to run jobs for up to 14 days (instead of the default 2 days)
- Further limits on Slurm Trackable RESources (TRES)

QOS	Partition	Allowed [L1] Account	Prio	${\tt GrpTRES}$	${\tt MaxTresPJ}$	MaxJobPU	Flags
besteffort	*	ALL	1			100	NoReserve
low	*	ALL (default for CRP/externals)	10			2	DenyOnLimit
normal	*	Default (UL,Projects,)	100			50	DenyOnLimit
long	*	UL,Projects,etc.	100	node=12	node=2	4	DenyOnLimit,PartitionTimeLimit
debug	interactive	ALL	150	node=8		2	DenyOnLimit
high	*	(restricted) UL, Projects, Industry	200			10	DenyOnLimit
urgent	*	(restricted) UL, Projects, Industry	1000			100	DenyOnLimit
admin	all	(restricted) sysadmins	1000				DenyOnLimit





# **ULHPC** Fairsharing 2.0

- Three fairsharing algorithms implemented in Slurm:

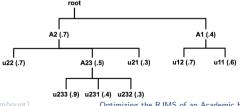
  - $\,\hookrightarrow\,$  Thourough evaluation of all 3 fairshare algorithms (Python-based simulator)



# **ULHPC** Fairsharing 2.0

- Three fairsharing algorithms implemented in Slurm:

  - → Thourough evaluation of all 3 fairshare algorithms (Python-based simulator)
- In Fair Tree, all users from higher priority account receive higher fair share factor
  - $\hookrightarrow$  ... when compared to all users from a lower priority account
  - → Done with rooted plane tree (rooted ordered tree)
    - $\checkmark$  logically created then sorted by fairshare level with highest fairshare values on the left
    - √ tree is then visited in a depth-first traversal way







# **ULHPC Fairsharing 2.0**

- Three fairsharing algorithms implemented in Slurm:

  - → Thourough evaluation of all 3 fairshare algorithms (Python-based simulator)
- New configuration with Multifactor Priority Plugin and Fair tree algorithm
  - → efficiency: new jobs are immediately assigned a priority.

  - → YET quite sensitive to the raw shares associated to each user account
- Necessity to deeply restructure associations & shares in the accounting DB
  - → formalize consistent rules to attribute raw shares.
  - → give novel incentives to good practices
    - √ take into account monetary contribution to increase priority
    - √ tribute to past efficient usage





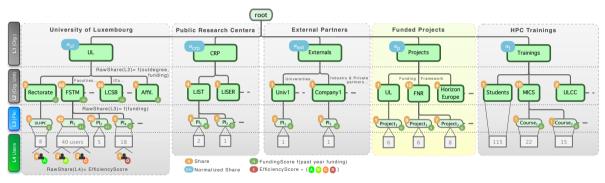
# **Account Hierarchy 2.0**

- Accounting records re-organized as a hierarchical tree (3 layers  $L_{1,2,3}$  + leafs)
  - - $\checkmark$  guarantee 85% of the shares for core UL activities
  - $\hookrightarrow$  L2: Organizational Unit (Faculty, ICs, External partner, Funding program...)
    - √ raw share depends on outdegree and funding score
  - - √ raw share depends on funding score (different weight)
    - √ eventually restricted only to projects and courses
  - $\hookrightarrow$  L4: End User (ULHPC login)
    - √ Raw share based on efficiency score



#### Toward Improved User Job Management through a Novel RJMS configuration

# **Account Hierarchy 2.0**







# Funding Score (L2/L3)

- Associated with account A belonging to level L in the hierarchy

  - $\hookrightarrow$  depreciation based on contribution type, weighted by level threshold  $eta_L$

$$FundingScore_L(A) = \left\lfloor \beta_L \frac{Investment_A(Year - 1)}{\#months} \right\rfloor$$





# Funding Score (L2/L3)

- Associated with account A belonging to level L in the hierarchy
  - $\hookrightarrow$  yearly updated at the beginning of the year
  - $\hookrightarrow$  depreciation based on contribution type, weighted by level threshold  $eta_{L}$

$$FundingScore_L(A) = \left\lfloor \beta_L \frac{Investment_A(Year - 1)}{\#months} \right\rfloor$$

- Ex1: Exceptional contribution of 120K€ performed in 2021 by a faculty (L2 account A)

  - $\hookrightarrow$  funding score in 2022:  $\left|\beta_{L_2} \frac{120000}{12}\right| = \left|\beta_{L_2} \times 10000\right|$ .
- Ex2: let P be a project granted in 2021 to start in 2022 for a duration of 36 months
  - budget: 27K€ allocated for HPC costs
  - $\hookrightarrow$  funding score for the years 2022, 2023 and 2024:  $\left\lfloor \beta_{L_3} \frac{27000}{36} \right\rfloor = \left\lfloor \beta_{L_3} \times 750 \right\rfloor$





# Efficiency Score (L4)

- Updated every year based on past jobs efficiency.
  - $\hookrightarrow$  Similar notion of "nutri-score': A(very good 3), B (good: 2), C (bad, 1), D(very bad 0)
- Proposed Metric for user *U*: Average Wall-time Accuracy (WRA) (higher the better)
  - $\hookrightarrow$  Defined for a given time period (past year)

$$egin{aligned} S_{ ext{efficiency}}(\textit{U}, \textit{Year}) &= \mathsf{WRA}(\textit{U}, \textit{Year}) \ &= rac{1}{\textit{N}} \sum_{\textit{JobID}} rac{T_{ ext{elapsed}}(\textit{JobID})}{T_{ ext{asked}}(\textit{JobID})} \end{aligned}$$





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• U raw share:  $1 + S_{\text{efficiency}}(U, Year)$ 

Score	Avg. WR
A (3) very good B (2) good C (1) bad D (0) very bad	$S_{ m efficiency} \geq 75\%$ $50\% \leq S_{ m efficiency} < 75\%$ $25\% \leq S_{ m efficiency} < 50\%$ $S_{ m efficiency} < 25\%$





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• WIP: integrate other efficiency metrics (CPU, mem, GPU efficiency)





# **Job Accounting and Billing**

- Utilization of the University computational resources is charged in Service Unit (SU)
  - $\,\hookrightarrow\,$  1 SU  $\simeq$  1 hour on 1 physical processor core on regular computing node
  - Usage charged 0,03€ per SU (VAT excluded) (external partners, funded projects etc.)
- A **job** is **characterized\*** (and thus billed) according to the following elements:
  - $\hookrightarrow$   $T_{\text{exec}}$ : Execution time (in hours)
  - $\rightarrow$   $N_{\text{Nodes}}$ : number of computing nodes, and **per node**:
    - ✓ N<sub>cores</sub>: number of CPU cores allocated per node
    - √ Mem: memory size allocated per node, in GB
    - $\sqrt{N_{\rm gpus}}$ : number of GPU allocated per node
  - $\hookrightarrow$  associated weighted factors  $\alpha_{\it cpu}, \alpha_{\it mem}, \alpha_{\it GPU}$  defined as TRESBillingWeight in Slurm
    - $\checkmark$  account for consumed resources other than just CPUs, taken into account in fairshare factor
    - $\sqrt{\alpha_{cpu}}$ : normalized relative perf. of CPU processor core (reference: skylake 73,6 GFlops/core)
    - $\checkmark$   $\alpha_{mem}$ : inverse of the average available memory size per core
    - $\checkmark$   $\alpha_{GPU}$ : weight per GPU accelerator





# **Job Accounting and Billing**

#### Number of SU associated to a job

$$N_{\mathsf{Nodes}} \times [\alpha_{\mathit{cpu}} \times N_{\mathsf{cores}} + \alpha_{\mathit{mem}} \times \mathit{Mem} + \alpha_{\mathit{gpu}} \times N_{\mathsf{gpus}}] \times T_{\mathsf{exec}}$$

#### Current billing weights:

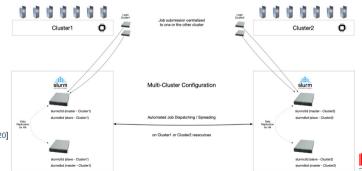
Cluster	Node Type	Partition	#Cores/node	CPU	$lpha_{ m cpu}$	$lpha_{ ext{mem}}$	$lpha_{ extsf{GPU}}$
Iris, Aion	Regular	interactive	28/128	n/a	0	0	0
ris	Regular	batch	28	broadwell	0.522	$\frac{1}{4} = 0,25$	0
ris	Regular	batch	28	skylake	1.0	$\frac{1}{4} = 0, 25$	0
ris	GPU	gpu	28	skylake	1.0	$\frac{1}{27}$	50
ris	Large-Mem	bigmem	112	skylake	1.0	<u>1'</u>	0
Aion	Regular	batch	128	ерус	0,57	$\frac{2}{1.75}$	0





# Federated vs. Multi-Cluster Scheduling Strategy

- Two possible approaches to integrating aion into the existing RJMS configuration
  - → Multi-Cluster: allow to submit jobs across each clusters sbatch -M {iris,aion} [...]

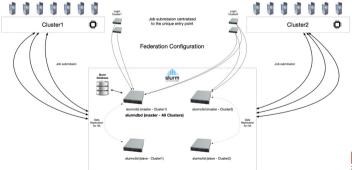


[Source: JOG, Oct 2020]



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  - → Federation: P2P replication of job submitted locally, distributed schedule attempt
    - √ unified job ID



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# Federated vs. Multi-Cluster Scheduling Strategy

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  - → Multi-Cluster: allow to submit jobs across each clusters sbatch -M {iris,aion} [...]
  - $\hookrightarrow$  Federation: P2P replication of job submitted locally, distributed schedule attempt  $\checkmark$  unified job ID

#### Selected Setup: Hybrid Multi-Cluster Scheduling Strategy

- Redundant master/slave RJMS controllers associated to each cluster
- Shared (like for federations), the same Slurm database service (slurmdb)





#### Impact Analysis of the Updated Models and Policies through Workload Performance Evaluation

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# Performance Evaluation and Experimental Setup

- Impact of the updated RJMS policy and setup particularly hard to qualify
  - → Approach 1: Slurm simulators of workloads execution. Ex: BSC slurm\_simulator [PMBS18]
    - ✓ **YET** old/obsolete version of Slurm featured (17.x)
    - √ novel proposed contributions absent and thus to implement...

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  - → Approach 2: Slurm Replay Engines. Ex: CSCS slurm-replay [SC18]
    - $\checkmark$  suffers from the same pittfalls: old/obsolete version of Slurm featured (18.x)

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[SC18] M. Martinasso & al. "RM-Replay: A High-Fidelity Tuning, Optimization and Exploration Tool for Resource Management", in Proc. of the Int. Conf. for High Performance Computing, Networking, Storage, and Analysis (SC18) (2018)



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### Our chance (Approach 3)

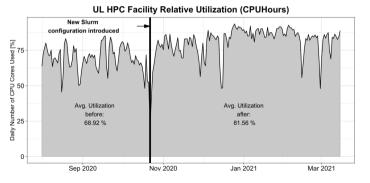
- All the modifications/proposals applied at once within our production systems!
  - $\rightarrow$  on Oct 22, 2020 during a maintenance session.
  - → 1 supercomputer (iris) has thus seen both configs for a sufficient amount of time
    - √ workload analysis able to reasonably capture impact of the changes on RJMS perf.





## Impact on the Total Utilization

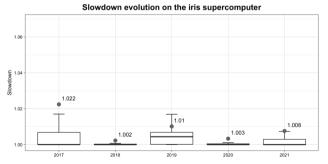
- Impact of the updated Slurm configuration on the ULHPC relative utilization:
  - $\hookrightarrow$  aggregates traces from several months of uninterrupted HPC services
  - $\hookrightarrow$  daily number of CPU cores used increased by 12.64%:  $\simeq$  81.56% over 6 months





## Impact on the Job Slowdown

- Slowdown: traditional performance metrics for job scheduling strategies [Feitelson15] slowdown =  $\frac{T_w + T_r}{T}$ 
  - RJMS response time normalized by the runtime:
  - $\rightarrow$  average yearly slowdown very close to the optimal value (1)

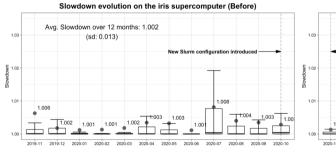


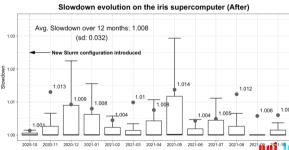
[Feitelson15] D. G. Feitelson, "Workload Modeling for Computer Systems Performance Evaluation, 1st ed. USA: Cambridge University Press (2015



## Impact on the Job Slowdown

- More detailed analysis to better capture the impact of the updated RJMS config.
  - → Monthly slowdown for all iris jobs run for 1 year period before and after the changes
  - → Fairly negligible impact: observed avg. job slowdown only increased by 0.59%



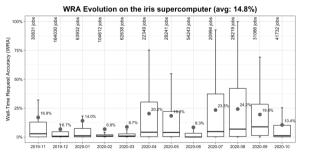




## Impact on the Average Wall-time Request Accuracy

Evaluation covering 1 year period before

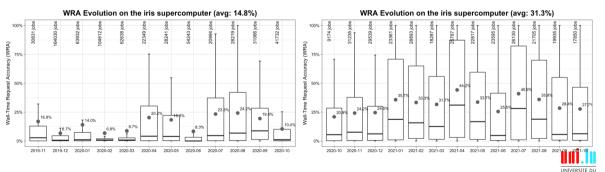
proposed configuration change





## Impact on the Average Wall-time Request Accuracy

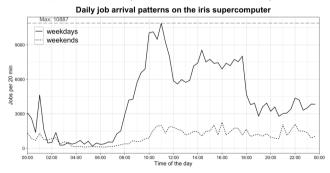
- Evaluation covering 1 year period before and after proposed configuration change
- Avg. WRA for the processed jobs was increased by 110.81%
  - $\hookrightarrow$  moving from 14.8% on average to 31.3%





## **Current Daily Arrival Pattern**

- Proposed changes in production for 18 months and extensively used







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### **Conclusion**

- In this talk:
  - $\hookrightarrow$  Optimization of Slurm RJMS config when acquiring & integrating a new supercomputer
    - √ smooth integration within the existing HPC ecosystem
    - √ mitigation of the identified pittfalls of the initial configuration after 3y of production





### **Conclusion**

#### In this talk:

- → Optimization of Slurm RJMS config when acquiring & integrating a new supercomputer
  - ✓ smooth integration within the existing HPC ecosystem
  - √ mitigation of the identified pittfalls of the initial configuration after 3y of production
- → Novel and flexible setup, adaptable to other HPC centers
  - ✓ new scheme defining partitions, QOS queues, priorities & their resource limits
  - $\checkmark$  fairsharing mechanism updated, with **incentives** to good practices & monetary contribution
  - √ rededigned account hierarchy, hybrid multi-cluster scheduling strategy



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- → New cost model and updated policies in production for 18 months
- → Performance evaluation from real workload traces
- Perspectives and Future directions
  - $\hookrightarrow$  smooth integration with Euro-HPC infrastructures
    - √ transparently outsource Research Computing/data analytic workflows to Tier-0 systems
  - → model & automatically offload from RJMS some of the less-demanding jobs
    - √ target dynamically allocated virtual cloud resources (burst instances)





#### Thank you for your attention...



## **Questions?**

Sebastien Varrette, Emmanuel Kieffer and Frederic Pinel Optimizing the Resource and Job Management System of an Academic HPC & Research Computing Facility – IEEE ISPDC 2022 University of Luxembourg. Belval Campus

Maison du Nombre, 4th floor 2, avenue de l'Université L-4365 Esch-sur-Alzette

mail: firsname.lastname@uni.lu

#### High Performance Computing @ Uni.lu

mail: hpc@uni.lu

Introduction: Context and Motivations

Toward Improved User Job Management through a Novel RJMS configuration

Impact Analysis of the Updated Models and Policies through

4 Conclusion & Perspectives

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#### ULHPC Technical Docs

hpc-docs.uni.lu



# Appendix / Backup Slides





### **ULHPC Job Prioritization Factors**

- Age: length of time a job has been waiting (PD state) in the queue
- Fairshare: difference between the portion of the computing resource that has been promised and the amount of resources that has been consumed
- Partition: factor associated with each node partition
  - $\hookrightarrow$  Ex: privilege interactive over batch
- QOS A factor associated with each Quality Of Service (low → urgent)

```
Job_priority =
PriorityWeightAge * age_factor +
PriorityWeightFairshare * fair-share_factor+
PriorityWeightPartition * partition_factor +
PriorityWeightQOS * QOS_factor +
- nice_factor
```





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```
# Current weights on ULHPC platform

$ sprio -w # --format "%8i %5A %9F %9P %Q"

JOBID AGE FAIRSHARE PARTITION QOS

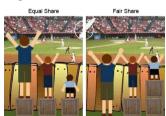
Weights 2000 3000 10000 10000
```





## **ULHPC Fairsharing 2.0**

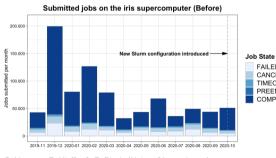
- Fairsharing: way of ensuring that users get their appropriate portion of a system
  - → Share: portion of the system users have been granted.
  - Usage: amount of the system users have actually used.
  - → Fairshare score: value the system calculates based off of user's usage.
    - √ difference between the portion of the computing resource that has been promised and the amount of resources that has been consumed
  - **Priority score**: priority assigned based off of the user's fairshare score.

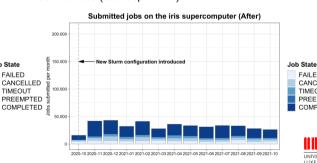




## Impact on the Daily submitted Jobs

- Obviously depends on the usage pattern of the platform
  - $\hookrightarrow$  Avg. #monthly submitted jobs: 71102  $\rightarrow$  35384
  - → Submission pattern changed and hard to conclude
    - COMPLETED job decreased...
    - Yet roughly eq. to CANCELLED+TIMEOUT increase (user-dependent)





Job State

FAILED

TIMEQUIT

COMPLETED

PREEMPTED

CANCELLED

Quantity Before

73.3%

0.12%

13.1%

10.4%

3.06%

Quantity After

60.3%

0.02%

15.2%

14.5%

9.92%

Difference

-13%

-0.1%

+2.1%

+4.1% +6.86%

FAIL ED

CANCELLED

PREEMPTED

COMPLETED

TIMEOUT